

# Provable Matrix Completion using Alternating Minimization

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# Alternating Minimization (AltMin)

## General Algorithm

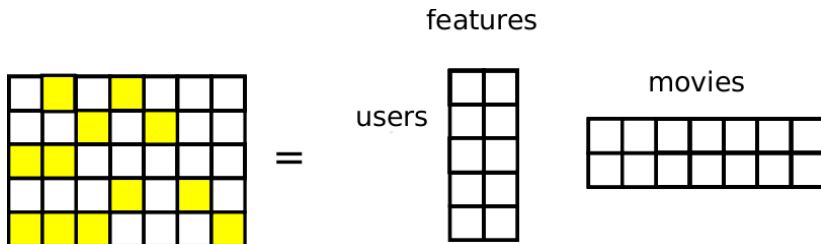
To minimize  $f(X)$  over rank- $k$  matrices  $X$ , repeat the following:

- fix  $U$  and minimize  $f(UV^\dagger)$  over  $V$
- fix  $V$  and minimize  $f(UV^\dagger)$  over  $U$

$$\begin{array}{ccc}
 X & = & U \quad V' \\
 \begin{array}{|c|c|c|c|c|c|c|c|} \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \end{array} & & \begin{array}{|c|c|c|c|c|c|} \hline \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square \\ \hline \end{array} \quad \begin{array}{|c|c|c|c|c|c|c|c|} \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \square & \square & \square & \square & \square & \square & \square & \square \\ \hline \end{array} \\
 m \times n & & m \times k & & k \times n
 \end{array}$$

- A popular Empirical approach to solve low rank matrix problems eg. matrix completion, clustering etc.
- **Challenge:** few theoretical guarantees

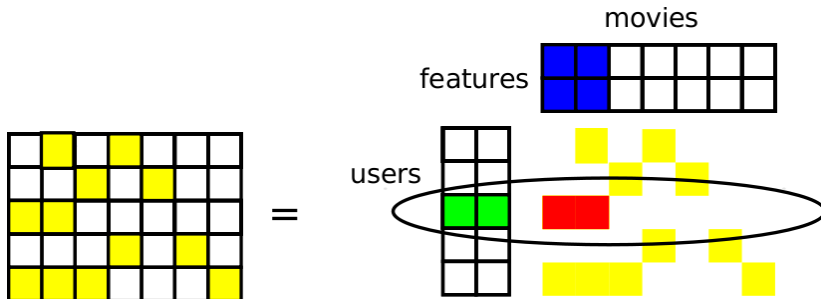
# Matrix Completion



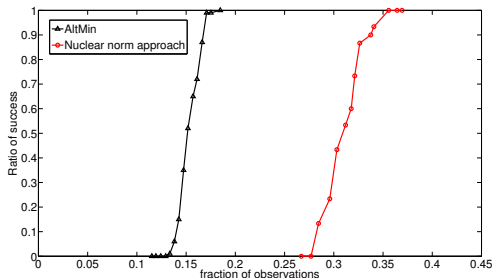
- Given some elements, fill in the rest
- Not possible in general; what if low rank?
- Metrics: Sample complexity and Computational complexity

# Matrix Completion via Alternating Minimization

$$\begin{aligned} \min \quad & \sum_{(i,j) \in \text{known set}} (M_{ij} - X_{ij})^2 \quad \text{s.t.} \quad \text{rank}(X) \leq k \\ = \min \quad & \sum_{(i,j) \in \text{known set}} (M_{ij} - U_i^\dagger V_j)^2 \quad \text{s.t.} \quad U \in \mathbb{R}^{m \times k}, V \in \mathbb{R}^{n \times k} \end{aligned}$$



# A Comparison



- Nuclear norm / Trace norm approach : convex relaxation.
- Empirically, AltMin has
  - similar sample complexity and
  - better computational complexity.

Challenge: AltMin formulation is non-convex.

# Our Results

- First theoretical guarantees for AltMin in any low rank setting
- We prove results for
  - matrix sensing
  - matrix completion

# Matrix Sensing

$$y = \mathcal{A}(X) = \begin{bmatrix} \langle A_1, X \rangle \\ \vdots \\ \langle A_d, X \rangle \end{bmatrix}$$

vector                      linear operator                      Low rank matrix

**Problem:** Given  $y$  and  $\mathcal{A}$ , recover  $X$ .

## Natural Algorithm (AltMinSense)

- 1 (Initialization)  $\hat{U}^0 \leftarrow$  top  $k$ -left s.v. of  $\sum y_i A_i$
- 2 In iteration  $t$ :
  - $\hat{V}^t \leftarrow \operatorname{argmin}_{V \in \mathbb{R}^{n \times k}} \left\| y - \mathcal{A}(\hat{U}^{t-1} V^\dagger) \right\|_2$
  - $\hat{U}^t \leftarrow \operatorname{argmin}_{U \in \mathbb{R}^{m \times k}} \left\| y - \mathcal{A}(U(\hat{V}^t)^\dagger) \right\|_2$

# Restricted Isometry Property (RIP)

Existing results require RIP assumptions.

## RIP [RFP10]

A linear operator  $\mathcal{A}(\cdot) : \mathbb{R}^{m \times n} \rightarrow \mathbb{R}^d$  satisfies  $k$ -RIP with  $\delta_k$ , if for all  $X \in \mathbb{R}^{m \times n}$  s.t.  $\text{rank}(X) \leq k$ , the following holds:

$$(1 - \delta_k) \|X\|_F^2 \leq \|\mathcal{A}(X)\|_2^2 \leq (1 + \delta_k) \|X\|_F^2.$$

- $\delta_k = 0 \Rightarrow$  Identity map
- $\delta_k = 1 \Rightarrow$  No information



# Existing Results

## Trace norm approach [RFP10]

$$\begin{array}{ll} \min \|y - \mathcal{A}(X)\|_2 & \rightarrow \min \|y - \mathcal{A}(X)\|_2 \\ \text{s.t. } \text{rank}(X) \leq k & \text{s.t. } \|X\|_* \leq \sqrt{k} \end{array}$$

- $\delta_{5k} < \frac{1}{10}$

## Singular Value Projection [JMD10]

- $\delta_{2k} < \frac{1}{3}$

## Drawback

- Need to compute many SVDs during execution - very slow in practice

# Our Results

## Theorem

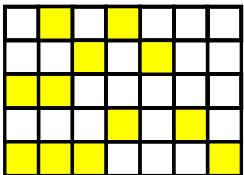
If  $\delta_{2k} < \left(\frac{\sigma_k^*}{\sigma_1^*}\right)^2 \frac{1}{100k}$ , then

$$\left\| M - \hat{U}^T (\hat{V}^T)^\dagger \right\|_F < \left(\frac{1}{2}\right)^T$$

## Remarks

- 1  $\delta_{2k}$  depends on the condition number unlike in existing work
  - modified algorithm:  $\delta_{2k} < \frac{1}{3200k^2}$
- 2 Linear convergence:  $\log \frac{1}{\epsilon}$  iterations for  $\epsilon$  error.

# Matrix Completion



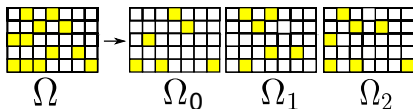
## Problem

Given elements in  $\Omega$ , find the low rank matrix  $M$ .

Analysis is harder

- $\Omega$  does not in general satisfy RIP.
- Dependence between iterates.

# Our Algorithm



- Divide  $\Omega$  into  $2T + 1$  subsets  $\Omega_0, \dots, \Omega_{2T}$  by uniform sampling.
- Use  $\Omega_i$  for the  $i^{\text{th}}$  iteration of AltMin.

## AltMinComplete

(Initialization)  $\hat{U}^0 \leftarrow$  top  $k$ -left s.v. of  $(M)_{\Omega_0}$

FOR  $t = 0, \dots, T - 1$

$$\hat{V}^{t+1} \leftarrow \operatorname{argmin}_{V \in \mathbb{R}^{n \times k}} \left\| \left( \hat{U}^t V^\dagger - M \right)_{\Omega_{t+1}} \right\|_F^2$$

$$\hat{U}^{t+1} \leftarrow \operatorname{argmin}_{U \in \mathbb{R}^{m \times k}} \left\| \left( U \left( \hat{V}^{t+1} \right)^\dagger - M \right)_{\Omega_{T+t+1}} \right\|_F^2$$

ENDFOR

- **Conjecture:** Do not need this partition.
- Same algorithm proposed and analyzed independently by [Kes12]

# A Hard Case

?	0			0		0
		0	0		0	
0	0			0		
0			0		0	
	0	0				0

# Incoherence

## Incoherence [CR09]

$M = U^* \Sigma^* (V^*)^\dagger$  is incoherent with parameter  $\mu$  if

- $\|u^{(i)}\|_2 \leq \frac{\mu\sqrt{k}}{\sqrt{m}} \forall i \in [m]$  and
- $\|v^{(j)}\|_2 \leq \frac{\mu\sqrt{k}}{\sqrt{n}} \forall j \in [n]$ .

coherent    incoherent

1	0
0	1
0	0
0	0
0	0
0	0
0	0

0	$-\frac{1}{\sqrt{n}}$
$\frac{1}{\sqrt{n}}$	$-\frac{1}{\sqrt{n}}$
$\frac{1}{\sqrt{n}}$	$-\frac{1}{\sqrt{n}}$
$\frac{1}{\sqrt{n}}$	0
$\frac{1}{\sqrt{n}}$	$\frac{1}{\sqrt{n}}$
$\frac{1}{\sqrt{n}}$	$\frac{1}{\sqrt{n}}$

# Existing Results

Existing results assume uniform sampling and incoherence of  $M$ .

Trace norm approach [CR09, CT09]

$$\min \sum_{(i,j) \in \Omega} (X_{ij} - M_{ij})^2 \quad \rightarrow \quad \min \sum_{(i,j) \in \Omega} (X_{ij} - M_{ij})^2$$

$$\text{s.t. rank}(X) \leq k \quad \text{s.t. } \|X\|_* \leq \sqrt{k}$$

- $O(kn \log n)$  observations
- Drawback: need many SVD calculations

OptSpace [KMO10]

- Opt. on Grassman manifold:  $O\left(f\left(\frac{\sigma_1^*}{\sigma_k^*}\right) kn \log n\right)$ 
  - rate of convergence not known

# Our Results

## Theorem

Let  $M$  be incoherent. If

$$\# \text{ measurements} > C \left( \frac{\sigma_1^*}{\sigma_k^*} \right)^6 k^7 n \log n \log \frac{1}{\epsilon},$$

then after  $T = O\left(\log \frac{1}{\epsilon}\right)$  iterations, we have:

$$\left\| M - \hat{U}^T (\hat{V}^T)^\dagger \right\|_F < \epsilon.$$

### Advantages:

- linear convergence :  $\log \frac{1}{\epsilon}$  vs  $\frac{1}{\sqrt{\epsilon}}$
- each iteration very fast
- low storage requirement

### Weakness: Dependence on

- condition number
- required accuracy
- $k$



# Comparison

	Sample comp. ( $d$ )	Comp. comp.
Our Results	$O\left(\left(\frac{\sigma_1^*}{\sigma_k^*}\right)^6 k^7 n \log n \log \frac{1}{\epsilon}\right)$	$O\left(dk^2 \log \frac{1}{\epsilon}\right)$
AltMin [Kes12]	$O\left(\left(\frac{\sigma_1^*}{\sigma_k^*}\right)^8 kn \log n \log \frac{1}{\epsilon}\right)$	$O\left(dk^2 \log \frac{1}{\epsilon}\right)$
Trace norm [CT09]	$O(kn \log n)$	$O\left(\frac{n^3}{\sqrt{\epsilon}}\right)$

# Main Idea of the Proof

- If  $\Omega =$  all elements, then AltMin becomes the well-known power method.
- In general, iterates take the form:

$$\widehat{V}^{t+1} = \underbrace{V^* \Sigma^* U^{*\dagger} U^t}_{\text{Power-method Update}} - \underbrace{F}_{\text{Error Term}} \text{ and}$$

$$\|F\|_2 \downarrow \text{ as } t \uparrow .$$

- Use RIP to show decay.
- Technical difficulty: Establishing incoherence of  $U^t$ .

# Summary and Beyond

## Summary

- First theoretical guarantees for AltMin in any low rank setting
- Results for
  - Matrix sensing
  - Matrix completion

## Further Directions

- Recent result for AltMin in Phase retrieval [NJS13]
- Theory for AltMin in clustering, sparse PCA, NMF etc.?

Thank you!

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